# Introduction to Topological Data Analysis - I 

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Motivation

[3D images (porous rocks)]

[Force fields in granular media]

[Nano-materials -
Li et al 2017]

Data often have topological / geometric structure

How many mountains?


How many mountains?


## Height is insufficient



Small bump on the side of a
Independent mountain
bigger mountain

## Prominence (Topography)

Local maximum: how low do you need to go before you can reach a higher maximum?


## Superlevelsets

$f: X \rightarrow \mathbb{R}$
$F_{t}=f^{-1}([t,+\infty))=\{x \in X, f(x) \geq t\}$
$F_{+\infty}=\emptyset, \quad F_{t} \subseteq F_{t^{\prime}}$ when $t \geq t^{\prime}, \quad F_{-\infty}=X$

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Smaller mountain merged into a bigger mountain: end the bar

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## Barcode

What it contains, for each local maximum

- height
- prominence


## Things it ignores

- position of the local maximum
- when a bar stops, which other bar it merges with
- what mountains are adjacent
- width of the mountains
- invariant by reparametrization

Why not the merge tree?


Why not the merge tree?


## Persistence diagram



Higher dimension


Higher dimension


## Higher dimension

Mathematical tool: homology
Defines "holes" of all dimensions. dim 0: connected components, dim 1: loops, etc.


2 connected components, 2 loops


1 connected component, 2 loops, 1 cavity

## Higher dimension

$f: \mathbb{R}^{2} \rightarrow \mathbb{R}$

Superlevelsets: sweep a horizontal plane
local maximum: new connected component
saddle point: merge 2 components, or create loop
local minimum: kill (fill) loop


One barcode / diagram per dimension (often drawn together in different colors)

Higher dimension


## Stability




## Stability




$$
\|f-g\|_{\infty}=\sup _{x \in X}\{|f(x)-g(x)|\}
$$

## Bottleneck distance

Partial matching, the rest matched with the diagonal
The worst pair defines the cost
Sup norm between points: $\max \left(\left|x-x^{\prime}\right|,\left|y-y^{\prime}\right|\right)$
Minimum over all matchings ( $\sim$ Wasserstein $W_{\infty}$ )
Can also define other distances $W_{p}$


## Stability Theorem

$$
d_{B}(\operatorname{Dgm}(f), \operatorname{Dgm}(g)) \leq\|f-g\|_{\infty}
$$

Dgm is 1-Lipschitz

## Independent (?) problem: point clouds

Input: point set $P$
Assumption: $P$ approaches some unknown ideal object
What can we do?

Strong reconstruction


## Strong reconstruction

"Connect the dots" homeomorphic reconstruction


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"Connect the dots" homeomorphic reconstruction Diffeomorphism?


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Diffeomorphism?
Requires a nice sampling

## Strong reconstruction

"Connect the dots" homeomorphic reconstruction Diffeomorphism?

Requires a nice sampling
Requires hypotheses on the model

Weaker reconstruction


## Weaker reconstruction

Thickened version of the object
Not homeomorphic to a circle
Same homotopy type


## Even weaker

Clustering
Mapper (graph) - next class
Persistent homology

## Topology of points?

Just $n$ connected components...

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Walk back, blur

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Walk back, blur
1 connected component, 1 loop

Choosing the scale


III defined problem
$\Longrightarrow$ Look at all scales!

Link with functions

$$
\begin{aligned}
& f: \mathbb{R}^{d} \rightarrow \mathbb{R} \\
& f(x)=\min _{p \in P}\|x-p\|_{2}
\end{aligned}
$$

Union of balls $=$ sublevelset of $f$

Persistence of offset filtration


Persistence of offset filtration


Persistence of offset filtration


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Persistence of offset filtration


## Persistence of offset filtration



Persistence of offset filtration


Persistence of offset filtration


Persistence of offset filtration


## Stability Theorem

$$
d_{B}\left(\operatorname{Dgm}\left(f_{P}\right), \operatorname{Dgm}\left(f_{Q}\right)\right) \leq\left\|f_{P}-f_{Q}\right\|_{\infty}=d_{H}(P, Q)
$$

$d_{H}$ : Hausdorff distance

$$
d_{H}(A, B)=\max \left\{\sup _{b \in B} d(b, A), \sup _{a \in A} d(a, B)\right\} \quad \text { where } d(b, A)=\inf _{a \in A} d(b, a)
$$



## How can we compute all that?

Computers need finite representations
Homology needs a notion of boundary
Hole $\sim$ gluing 2 regions with the same boundary

$\Longrightarrow$ Cell complexes (generalization of graphs)

Cubical


Simplicial


## Filtered cubical complex

The cells are cubes of all dimensions (including vertices, edges, squares)
Represent a region of $\mathbb{R}^{d}$ by a subset of cells (subcomplex)
Growing region $=$ sequence of subcomplexes $K_{0} \subseteq K_{1} \subseteq \ldots \subseteq K_{n}$


Each cell $\sigma$ has a filtration value $f(\sigma)$ : time of appearance
Subcomplex $K_{t}=\{\sigma, f(\sigma)<=t\}$
Homology needs boundaries: the faces of a cell of $K_{i}$ are also cells of $K_{i}$
$\sigma \subset \tau \Longrightarrow f(\sigma) \leq f(\tau)$
Can compute the persistence diagram of the sequence of subcomplexes

## Persistence algorithm in dim 0

Only uses a (filtered) graph
Insert vertices and edges one by one (filtration order)
Vertex: new connected component
Edge $a b$ : if $a$ and $b$ in separate components, kill the youngest, otherwise new loop (ignored for $\operatorname{dim} 0$ )


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Edge $a b$ : if $a$ and $b$ in separate components, kill the youngest, otherwise new loop (ignored for dim 0 )
Disjoint-set data structure (aka union-find): very fast, running time dominated by sorting edges
Dimension $p$ :

- Use cells of dimension $p$ and $p+1$
- Replace "separate components" with algebra (boundary not in the vector space generated by previous boundaries)
- Worst case $\Theta\left(n^{3}\right)$, in practice $O(n) \quad$ ( $n$ number of cells)
- Cohomology? Same diagram

Functions
Discretize the function: grid points


## Functions

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Extend to other cells: lower-star filtration
$f(\sigma)=$ maximum value at its vertices
Same persistence diagram as a piecewise linear interpolation


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level 0

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level 1

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level 2

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level 3

## Functions

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level 4

## Simplicial complex

Simplex: vertex, edge, triangle, tetrahedron, etc


Simplices have the minimal number of vertices for their dimension (no squares or pentagons)
Nice combinatorial intersection: $\sigma \cap \tau$ is a common face of $\sigma$ and $\tau$ or empty.
Represent a region by a subset of cells (subcomplex)


Growing region $=$ sequence of subcomplexes $K_{0} \subseteq K_{1} \subseteq \ldots \subseteq K_{n}$
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Can compute the persistence diagram of the sequence of subcomplexes


Nerve of a cover $\left\{U_{1}, \ldots, U_{n}\right\}$
One vertex $v_{i}$ per $U_{i}$
$\sigma=\left[v_{i_{0}}, \ldots v_{i_{k}}\right] \in K \Longleftrightarrow \bigcap_{j=0}^{k} U_{i_{j}}$
Abstract (not embedded)
Nerve Theorem: If all intersections $\bigcap_{i \in I \subseteq[1, n]} U_{i}$ are either empty or contractible, then the union $\bigcup_{i=1}^{n} U_{i}$ has the same homotopy type as the Nerve (simplicial complex).

Always true for convex objects (balls)


Persistent nerve of a growing cover $\left\{U_{1}^{t}, \ldots, U_{n}^{t}\right\}$
$U_{i}^{t} \subseteq U_{i}^{t^{\prime}}$ when $t \leq t^{\prime}$
$K^{t}$ : Nerve of $\left\{U_{1}^{t}, \ldots, U_{n}^{t}\right\}$
$K^{t} \subseteq K^{t^{\prime}}$ when $t \leq t^{\prime}$
Persistence diagram of $U^{t}=\bigcup_{i=1}^{n} U_{i}^{t}$
Persistence diagram of $K^{t}$


Persistent nerve theorem: If all intersections $\bigcap_{i \in I \subseteq[1, n]} U_{i}^{t}$ are either empty or contractible, then the two diagrams are the same.

## Čech filtration

Finite point set $P$, parameter $r$
$C_{r}(P)$ is the nerve of the union of the balls of radius $r$ centered on the points of $P$
The sequence of $C_{r}(P)$ when $r$ increases defines a filtered simplicial complex
$f(\sigma)=$ radius of Minimal Enclosing Ball of the vertices of $\sigma$


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1 persistence diagram for the growing union of balls when $r$ increases
1 persistence diagram for the sequence of Čech complexes when $r$ increases
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1 persistence diagram for the growing union of balls when $r$ increases
1 persistence diagram for the sequence of Čech complexes when $r$ increases
Persistent nerve theorem: those 2 diagrams are the same
Weaknesses:
Big, $2^{|P|}$ simplices if we do not limit it
Numerical, geometric computation to test the intersection of $k$ balls in $\mathbb{R}^{d}$

## Rips filtration

Finite point set $P$, parameter $r$

$$
\sigma \in R_{r}(P) \Longleftrightarrow \operatorname{diam}(\sigma) \leq 2 r
$$

Same graph as the Čech complex $C_{r}(P)$
Add simplices for cliques (complete subgraphs): purely combinatorial
$C_{r}(P) \subseteq R_{r}(P) \subseteq C_{2 r}(P)$
$\Longrightarrow$ persistence diagrams of Rips and Čech are at distance less than 2 in log-scale

Very flexible (arbitrary value on edges)
Drawbacks: not exact topology, and even bigger than Čech
Stability theorem still applies

## $\alpha$-complex

Čech: lot of redundancy in the cover when $r$ is large.
Idea: no need to keep growing in places already covered by other balls


## Choice: Rips or $\alpha$-complex?

Seldom compute the true Čech complex
Rips: easy, very flexible. Independent of any embedding. Have to limit dimension and edge length $\alpha$-complex: only defined in Euclidean $\mathbb{R}^{d}$, currently only efficient for small $d$, but super efficient there Several other alternatives, including sparse Rips, etc

## Noise (outliers)

Stability theorem: only handles small perturbations, one outlier can break everything
Idea: superlevelsets of a density estimator
Computable version: weighted version of Čech, Rips; penalize points in low density regions


## Crossover: time series as a point cloud

Time series: function $f: \mathbb{R} \rightarrow X$
If $X=\mathbb{R}$ use sublevelsets?
Alternative idea: forget time, see $f(\mathbb{R}) \subseteq X$ as a point set, compute Čech complex

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$$
X=\mathbb{R}^{2}
$$

$6 \cdot$

- 3
$1{ }^{\bullet}$
- 2


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Lose too much information? Enrich $f$ first, define $g=\left(f, \frac{\partial f}{\partial t}\right)$



$$
f(\mathbb{R})
$$

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In practice, from sequence $u_{n}$, define e.g. $v_{n}=\left(u_{n}, u_{n+2}, u_{n+7}\right)$
Inspired by Taken's theorem in dynamical systems

## Conclusion

Can be expensive to compute: aim for a smaller complex
Hardest part: deciding on what function to compute persistence
Next classes: how to use this in stat / ML

